TECHNICAL REPORT ON MACHINE LEARNING EXPERIMENTS FOR THE MÖBIUS FUNCTION

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ABSTRACT. Last week, I was at the [Mathematics and Machine Learning](https://cmsa.fas.harvard.edu/event/mml2024/) program at Harvard's Center of Mathematical Sciences and Applications. The underlying topic was on number theory and I've been studying various number theoretic problems from a machine learning perspective.

This is a technical report, including details related to actually running the code and analyzing the results.

CONTENTS

1. INTRODUCTION

I've been computing several experiments related to estimating the Mobius function $\mu(n)$. Previous machine learning experiments on studying $\mu(n)$ have used neural networks or classifiers. Francois Charton made an integer sequence to integer sequence transformer-based translator, [\[Int2Int\]](#page-23-1) (available at [Int2Int\)](https://github.com/f-charton/Int2Int/), and I thought it would be fun to see if this works any different.

Initially, I sought to get Int2Int to work. I describe aspects of that and how to run it in various ways here.

I'm splitting my description into two parts: a general report [\[DLD-General\]](#page-23-2) and a technical report. This is the technical report. This includes many details related to actually running and analyzing the code.

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2. Int2Int

Int2Int can be found at [Int2Int.](https://github.com/f-charton/Int2Int/) It is possible to run Int2Int using a CPU, but it's much slower and probably not worth trying.

Francois and Edgar Costa (and to a lesser extent, me) have tried ot make Int2Int as self-contained as possible. It's certainly easier now than it was a few weeks ago — it might be possible for the Reader to experiment with https://github.com/f-charton/Int2Int/ using only the README there.

2.1. **Training from data files.** By default, Int2Int expects to be able to generate valid inputs and outputs on the fly. We wanted to use and experiment with data that is nontrivial to compute (such as the Möbius function or data associated to elliptic curves). For that, we've added the ability to train from data files.

It would be fair to say that running Int2Int with default settings is so easy that the hardest part to get up and running is creating the data files. And this is only as hard as the data is to generate and store.

To train from data files, the file needs ot have a particular format. Recall that Int2Int fundamentall reads and outputs sequences of integers. Each integer is encoded as s ad ... a0, where s is either + or - and is the sign, and ad through a0 are the digits in a given base (which defaults to 1000). For example, the number 12345 is encoded as + 12 345.

An array of n integers is encoded as ν n $z1$... zn, where the n in ν n is the actual number. For example, the array (1,1234,1234567) is encoded as V3 + 1 + 1 234 + 1 234 567.

A datafile should have the input given as an array of the appropriate length, followed by a tab character \cdot , followed by the output. As an even more technical note, the output can be specified by a range of values instead of as an integer or integer array; this is useful with $\mu(n)$ since it can only take 3 values. This has to do with the symbol table that Int2Int uses, and the fact that it uses [cross-entropy loss](https://en.wikipedia.org/wiki/Cross-entropy) to measure performance.

A complete datafile could be the following.

 1 V5 + 1 + 2 + 1 + 3 + 4 + 5\t+ 1 $2 \quad V5 + 0 + 2 + 1 + 3 + 1 + 5 \tt{t}$

This data file has the spec int $[5]$: int. If we wanted more than a single int as output, we would have to use Vn appropriately. There is an additional datatype called range (with python-like semantics). In practice, if we know the output is a single constrained integer, there is a minor boost from using range instead of int.

2.2. **Datafile Generation Scripts.** I generated most of my datafiles using a script that closely looked like the following. This is a sagemath script. That is, it's mostly python, but it has inbuilt commands primes and moebius that I take for granted.

```
1 primes_100 = list(primes(542)) # generate list of 100 primes
\overline{2}3 def encode_integer(val, base=1000, digit_sep=" "):
\dot{ } if val == 0:
5 return '+ 0'
6 sgn = '+' if val >= 0 else '-'
\tau val = abs(val)
8 r = []9 while val > 0:
10 r.append(str(val % base))
11 val = val//base12 r.append(sgn)
13 r.reverse()
14 return digit_sep.join(r)
15
16 # Each line has an input, a tab, and an output.
17 def make_line(n):
18 return make_input(n) + "\t" + make_output(n) + "\n"
19
20 def make_input(n):
21 ret = [1]22 count = len(primes_100)
23 ret.append(f''V{2*count}'')24 for p in primes_100:
25 ret.append(encode_integer(n % p)) # feed in n mod p
26 ret.append(encode_integer(p)) # followed by p
27 return ' '.join(ret)
28
29 def make_output(n):
30 return str(moebius(n))
```
What's left is to determine what family of *n* to input. For performing regressionlike tasks, Francois noted that using a log-distribution tends to work best. This is more like a classification task as the output is one of 0, −1, or 1. Thus I generically uniformly sampled integers n up to some large bound like 10^{13} without repetition. To do this, I generate random integers in the range and check to make sure that I don't generate the same one twice.

```
1 import random
2
3 seen = set()
4 with open("mu_modp_and_p.txt", "w", encoding="utf8") as
   ,→ outfile:
5 while len(seen) < 10**7:
6 n = \text{random.random}(2, 10**13)7 if n in seen:
```
⁸ continue ⁹ seen.add(n) ¹⁰ outfile.write(make_line(n))

Note that this creates 10^7 lines, each having approximately 200 \cdot 3 ~ 1000 characters. The resulting file will be approximately 10GB. Adjust the parameters appropriately!

The slow part is computing $\mu(n)$ for random integers. Generating random numbers (including the 10^{-6} chance of hitting a previously seen number) and writing to the file is fast; computing $\mu(n)$ for a random 12 digit number can be slow-ish.

But in practice, the actual slow part is training the resulting ML model. I didn't work to optimize generation of μ at all.

I note that a sieve could generate all the Möbius values up to *N* at once. Then you could sample from these values in whatever way makes sense. Something along the following lines would work (and would remove the sagemath dependency).

```
1 def primes_up_to(X):
2^{\frac{1}{2}} ""
3 A basic implementation of Eratosthenes.
4 """"
5 arr = [True] * (X + 1)6 \arr[0] = \arr[1] = False7 primes = []
8 for p in range(X + 1):
9 if arr[p]: # is prime
10 primes.append(p)
11 for j in range(p*p, X + 1, p):
12 arr[j] = False
13 return primes
14
15
16 def mobius_up_to(X):
17 "Eratosthenes-like"
18 arr = [1] * (X + 1)19 arr[0] = 020 ps = primes_up_to(X)
21 for p in ps:
22 for j in range(p, X + 1, p):
23 arr[j] * = -124 for j in range(p*p, X + 1, p*p):
25 arr[j] = 026 return arr
```
2.3. **Making testing and training data.** I then make testing and training data.

```
1 import os
2 def shuffle_and_create(fname, ntrain=1900000, ntest=100000):
3 "Shuffle and create test and training files"
4 if not fname.endswith(".txt"):
5 raise ValueError("Incorrect filename assumption.")
6 name = fname [-4] # remove ".txt"
7 print("shuffling...")
8 os.system(f"shuf {name}.txt > {name}.shuf.txt")
9 print("making training data...")
10 os.system(f"head -n {ntrain} {name}.shuf.txt >
          {name}.txt.train")11 print("making testing data...")
12 os.system(f"tail -n {ntest} {name}.shuf.txt >
         {name}.txt.test")
13 print("done!")
```
2.4. **Running Int2Int.** It's now time to actually run the code. It's necessary to have a python with pytorch installed (not surprisingly) and to have Int2Int somewhere. But a generic run would look like

```
1 python ../Int2Int/train.py
2 --num_workers 0
3 --dump_path \tilde{\phantom{a}}/scratch
4 --exp_name dld_mu_modp_and_p_sqfree
5 --exp_id 1
6 --train_data ./mu_modp_and_p_sqfree.txt.train
7 --eval_data ./mu_modp_and_p_sqfree.txt.test
8 --local_gpu 1
9 --epoch_size 250000
10 --operation data
11 -- data_types "int[200]: range(-1,2)"
12 --optimizer "adam, 1r=0.00025"
```
This was one of the commands I used when using $(n \mod p, p)$ for the first 100 primes (giving 200 inputs total) and output just $\mu(n)$ (one int in a prescribed range).

Most of these are straightforward. The --optimizer command is deceptively useful, largely because changing the initial learning rate can have large impacts on the overall performance.

When run in this way, it's almost certain that you'll need to manually stop the experiment before it has a complete run. This is because the default number of epochs to train through is very large. In practice, it's a good idea to sometimes look at the outputs or parse the logs and to see how the behavior is going.

2.5. **Log parsing and graph creation.** For log parsing and graph creation, I used code largely written by someone else (maybe Edgar Costa). This is a pile of code, but it's just parsing the pickled logs from a set of experiments. Log writing and parsing always takes piles of not-very-hard code. This is no exception.

The relevant code is in the appendix.

3. Representation

I thought using residues mod several primes was a good strategy. Other experiments have shown^{[1](#page-5-2)} that the base in which numbers are expressed can be very important.

Something like base 100 or base 1000 would allow for almost immediate recognition that $\mu(n) = 0$ if $25 | n$ or if $4 | n$, as these congruence classes are fixed. I'm more interested in what other sorts of mathematical structures the machine can learn.

In this case, I represented each number in base 1000, but almost never needed to use any number larger than 1000 (as the 100th prime is 541). The Chinese remainder theorem shows that this allows representation for every integer up to approximately $10^{219.67}$. This is large enough to be interesting.

```
1 # pure python - uses primes_up_to defined above
2 import math
3 from functools import reduce
4
5 primes_100 = primes_up_to(1000)[:100]
6 print(primes_100[-1])
7 # 541
8
9 modulus = reduce(lambda x, y: x*y, primes_100, 1)
10 print(modulus)
11 # [...enormous...]12 print(math.log(modulus)/math.log(10))
13 # 219.67...
```
If $n < 10^{219.67}$, then *n* is uniquely determined by its residues mod *p* for the first 100 primes *p*.

4. GUESSING $\mu(n)$ from *n* mod *p* without using CRT

One of the questions that came up was the following **mathematical** (not programmatic) question.

How would you guess whether *n* is squarefree or not given *n* mod *p* for lots of primes *p*?

¹See *Learning the greatest common divisor: explaining transformer predictions* by François Charton. François also extracted Int2Int from the models used in this paper, more or less. Thank you François.

One way would be to perform the Chinese remainder theorem, reconstruct *n*, and then actually check. There is no known polynomial-time algorithm to check if an integer is squarefree, so this approach is generically slow.

The "default" algorithm would be to note that about 60.79 percent of numbers are squarefree. So guessing squarefree all the time would be right just over 60 percent of the time. I want any algorithm that does better.

The Dirichlet series for squarefree numbers that are divisible by a fixed prime *q* is

$$
\frac{1}{q^s} \prod_{\substack{p\\p \neq q}} \left(1 + \frac{1}{p^s} \right) = \frac{1}{q^s} \frac{(1 - 1/q^s)}{(1 - 1/q^{2s})} \frac{\zeta(s)}{\zeta(2s)},\tag{1}
$$

and the series for squarefree numbers that aren't divisible by a fixed prime *q* is the same, but without q^{-s} . Thus the percentage of integers that are squarefree and divisible by *q* or not divisible by *q* are, respectively,

$$
\frac{1}{q+1}\frac{6}{\pi^2} \quad \text{and} \quad \frac{q}{q+1}\frac{6}{\pi^2}.
$$
 (2)

A simple application of conditional probability shows that

$$
P(\text{sqfree}|\text{q-even}) = \frac{P(\text{sqfree and q-even})}{P(\text{q-even})} = \frac{q}{q+1} \frac{6}{\pi^2}
$$

$$
P(\text{sqfree}|\text{q-odd}) = \frac{P(\text{sqfree and q-odd})}{P(\text{q-odd})} = \frac{q^2}{q^2 - 1} \frac{6}{\pi^2}.
$$

I use the adhoc shorthand q -even to mean divisible by q , and q -odd to mean not divisible by *q*.

Let's quickly experimentally verify this. We make squarefree numbers with yet another Eratosthenes-type sieve.

```
1 def squarefree_up_to(X):
2 """"
3 Eratosthenes-like.
4 """"
5 arr = [True] * (X + 1)6 arr[0] = False7 ps = primes_up_to(int(X**.5) + 1)
8 for p in ps:
9 for j in range(p*p, X + 1, p*p):
10 arr[j] = False
11 ret = \begin{bmatrix} \end{bmatrix}12 for i in range(X + 1):
13 if arr[i]:
14 ret.append(i)
15 return ret
16
17
18 sfree = squarefree_up_to(10_000_000)
```

```
19 print(len(sfree)/10_000_000)
20 # 0.6079291
21
22 import math
23 print(6./math,pi**2)24 # 0.6079271018540267
```
As an aside, I note that this converges very quickly. Look at how close that is! One useless application of the Riemann Hypothesis is that is would guarantee how quickly the density of the number of squarefree numbers up to *X* would converge to $6/\pi^2$.

```
1 def ratio_sqfree_with(filterfunc):
2 return sum(1 for n in sfree if filterfunc(n))/len(sfree)
3
4 def is even(x):
5 return 1 if x \frac{9}{2} = 0 else 0
6 def is_odd(x):
7 return 1 if x \frac{9}{2} == 1 else 0
8
9 # even and sqfree
10 ratio_sqfree_with(is_even)
11 # 0.3333309756022536
12
13 # odd and sqfree
14 ratio_sqfree_with(is_odd)
15 # 0.6666690243977463
16
17 def is_3even(x):
18 return 1 if x \frac{9}{6} 3 == 0 else 0
19 def is 3odd(x):
20 return not is_3even(x)
21
22 ratio_sqfree_with(is_3even)
23 # 0.24999839619455624
24
25 ratio_sqfree_with(is_3odd)
26 # 0.7500016038054438
```
This agrees with the claim above that $1/(q + 1)$ of squarefree numbers are divisible by the prime *q*, and $q/(q+1)$ are not. The converse probabilities follow from basic probability, but to make sure:

```
1 sqfree_set = set(sfree) # for quick inclusion checking
2
3 def prob_sqfree_given(filterfunc):
4 sqfree_count = 0
5 total count = 0
```

```
6 for n in (x for x in range(10_000_000) if filterfunc(x)):
7 total_count += 1
8 if n in sqfree_set:
9 sqfree_count += 1
10 if total_count == 0:
11 return 0.0
12 return sqfree_count / total_count
13
14 # P(sqfree | divis by 2)
15 prob_sqfree_given(is_even)
16 # 0.4052832
17
18 # P(sqfree | not divis by 2)
19 prob_sqfree_given(is_odd)
20 # 0.810575
21
22 # P(sqfree | divis by 3)
23 prob_sqfree_given(is_3even)
24 # 0.45594380881123825
25 3/4 * 6/math.pi**2
26 # 0.45594532639052
27
28 # P(sqfree | not divis by 3)29 prob_sqfree_given(is_3odd)
30 # 0.6839217683921769
31 9/8 * 6/math.pi**2
32 # 0.68391798958578
```
These are very close to the theoretical computations above — again, it turns out that convergence is very quick.

4.1. **Compound Probabilities.** We'll compute joint probabilities theoretically in a moment. But we'll also experimentally find them.

Let's look at the probability using the small-prime strategy for the primes 2,3,5: if *n* is divisible by one of these, guess that *n* is not squarefree; otherwise guess that *n* is squarefree.

```
1 def not_divis_by_small_prime(n):
2 for p in (2, 3, 5):
3 if n \nmid p == 0:
4 return False
5 return True
6
7 A = prob_sqfree_given(not_divis_by_small_prime)
8 print(A)
9 # 0.9498902374725594
10
```

```
10 DAVID LOWRY-DUDA LAST UPDATED: 2024.10.19
11 def prob_notsqfree_given(filterfunc):
12 notsqfree_count = 0
13 total_count = 0
14 for n in (x for x in range(10_000_000) if filterfunc(x)):
15 total_count += 1
16 if n not in sqfree_set:
17 notsqfree_count += 1
18 if total_count == 0:
19 return 0
20 return notsqfree_count / total_count
21
22
23 def divis_by_small_prime(n):
24 return not not_divis_by_small_prime(n)
25
26 B = prob_notsqfree_given(divis_by_small_prime)
27 print(B)28 # 0.5164203621436034
```
The density of numbers not divisible by 2,3, or 5 is $(1-1/2)(1-1/3)(1-1/5) \approx$ 0.2666. Thus 0.2666 of the time, *n* isn't divisible by 2 or 3 or 5 and we would guess that *n* is squarefree; this is correct about 0.9498 of the time. And the 0.7333 of the time when *n* is divisible by at least one of 2 or 3 or 5, we guess that *n* is not squarefree; this is correct 0.5164 of the time.

In total, we expect that this strategy is correct with density

 $0.2666 \cdot 0.9498 + 0.7333 \cdot 0.5164 \approx 0.6318$.

Let's check:

```
1 not_divis_prob = (1 - 1/2)*(1 - 1/3)*(1 - 1/5)2 corr = not_divis_prob * A + (1 - not_divis_prob) * B
3 print(corr)
```

```
4 # 0.6320123288979917
```
If you look, you'll see that this does better than the naive guess (always guess squarefree) but is worse than guessing based only on mod 2 data. This is because we're ignoring all of the various cross-correlations. Clearly incorporating cross-correlations can never do worse than only using the mod 2 data.

Suppose we look instead at all the probabilities for all 2^{ℓ} possibilities of n being divisible or not by the first *ℓ* primes. Here, I use the first 4 primes, and the strategy is simple: compute whether it is more likely for *n* to be squarefree or not given each divisibility pattern, and guess that one.

```
1 def binary_to_prime_sets(n, length=4):
2 assert length <= 25
3 b = bin(n)[2:]4 b = "0" * (length - len(b)) + b5 is divis = []
```

```
6 \qquad \text{not\_divis = []}7 ps = primes_up_to(100)[:length]
8 for l, p in zip(b, ps):
9 if 1 == "1":
10 is_divis.append(p)
11 else:
12 not_divis.append(p)
13 return is_divis, not_divis
14
15 def divis_rules(is_divis, not_divis):
16 def filterfunc(n):
17 for p in is_divis:
18 if n % p != 0:
19 return False
20 for p in not_divis:
21 if n \gamma p == 0:
22 return False
23 return True
24 return filterfunc
25
26 def density_given(is_divis, not_divis):
27 filterfunc = divis_rules(is_divis, not_divis)
28 count = 029 for n in (x for x in range(10_000_000) if filterfunc(x)):
30 count += 1
31 return count/10_000_000
32
33 correct = 034 exp = 435 for n in range(2**exp):
36 is_divis, not_divis = binary_to_prime_sets(n, length=exp)
37 ff = divis_rules(is_divis, not_divis)
38 psqfree = prob_sqfree_given(ff)
39 density = density_given(is_divis, not_divis)
40 prob = max(psqrt{size}, 1 - psqfree)41 correct += density * prob
42 print(
43 is_divis, not_divis, n,
44 psqfree, density, prob, density * prob, correct
45 ) # my own diagnostics
46 print(correct)
47 # 0.7031860000000001
```
Remarkably this almost no better than just 2 alone! Before performing this computation, I had assumed that it would be notably better. Instead, it's close enough that it might actually be the same as using 2 alone, combined with numerical imprecision.

With this set up, we can compute the theoretical probabilities instead of using experimentally determined probabilities.

4.2. **Actual computation.** Let $\{p_1, \ldots, p_N\}$ and $\{q_1, \ldots, q_D\}$ denote two disjoint sets of primes. We want to compute the density of squarefree numbers that are divisible by each of the p_i and not divisible by any of the q_j . Each of these local conditions are independent; the overall density is the product of the local densities as described in \sim [\(1\)](#page-6-0) and \sim [\(2\)](#page-6-1). That is, the density of integers divisible by the p_i and not divisible by the q_j is

$$
\prod_{p_i}\left(\frac{1}{p_i+1}\right)\prod_{q_j}\left(\frac{q_j}{q_j+1}\right)\frac{6}{\pi^2}.
$$

Recall the chain rule from probability, that says

$$
P\left(\bigcap_{i=1}^k E_i\right) = P\left(E_1 | \bigcap_{i=2}^k E_i\right) = P\left(E_1 | \bigcap_{i=2}^k E_i\right) P\left(\bigcap_{i=2}^k E_i\right),
$$

(and which could chain further). I write P (sqfree, $p_1, p_2, \widehat{q_1}, \widehat{q_2}$) to mean the probability that a number is squarefree, divisible by p_1 and p_2 , and not divisible by *q*¹ or *q*² (with obvious notational generalization). Then

$$
P(\text{sqfree}|p_1,\ldots,p_N,\widehat{q_1},\ldots,\widehat{q_D})=\frac{P(\text{sqfree},p_1,\ldots,p_N,\widehat{q_1},\ldots,\widehat{q_D})}{P(p_1,\ldots,p_N,\widehat{q_1},\ldots,\widehat{q_D})}.
$$

Divisibility by different primes are independent, so this simplifies to

$$
P(\text{sqfree}|p_1,\ldots,p_N,\widehat{q_1},\ldots,\widehat{q_D})=\frac{P(\text{sqfree},p_1,\ldots,p_N,\widehat{q_1},\ldots,\widehat{q_D})}{P(p_1)\cdots P(p_N)P(\widehat{q_1})\cdots P(\widehat{q_D})}.
$$

We also have that $P(p) = 1/p$ and $P(\hat{q}) = (q-1)/q$.

Altogether, we compute that

$$
P(\text{sqfree}|p_1,\ldots,p_N,\widehat{q_1},\ldots,\widehat{q_D})=\prod_{p_i}\left(\frac{p_i}{p_i+1}\right)\prod_{q_j}\left(\frac{q_j^2}{q_j^2-1}\right)\frac{6}{\pi^2}.
$$

Note that this generalizes the previous probabilities and is generically straightforward.

Let's quickly check by computing P (sqfree|2,3) and P (sqfree|2,3):

$$
P(\text{sqfree}|2,3) = \frac{1}{2} \frac{6}{\pi^2} \approx 0.3039,
$$

\n
$$
P(\text{sqfree}|2,3) = \frac{3}{4} \frac{6}{\pi^2} \approx 0.4559.
$$

\n1 $\text{divis_by_2_and_3} = \text{divis_rules}([2, 3], [1])$
\n2 $\text{print}(\text{prob_sqfree_given}(\text{divis_by_2_and_3}))$
\n3 $\# 0.30395813920837217$
\n4
\n5 $\text{divis_by_2_not_3} = \text{divis_rules}([2], [3])$

```
6 print(prob_sqfree_given(divis_by_2_not_3))
```
0.45594574559457457

Let's now compute the density of the following strategy being correct:

- 1. Fix a set of primes *P*.
- 2. For each partition of *P* into two disjoint sets of primes $\{p_i\}$ and $\{q_i\}$:
- 3. Compute P (sqfree| p_1, \ldots, p_N , $\widehat{q_1}, \ldots, \widehat{q_D}$).
- 4. For integers satisfying this set of prime divisibility rules, guess "squarefree" if this probability is larger than 0.5; otherwise guess "not squarefree".

```
1 def prob_sqfree_theoretical(ps, qs):
```

```
2 ret = 6/math.pi**23 for p in ps:
4 ret *=(p / (p + 1))5 for q in qs:
6 ret *= (q * q / (q * q - 1))7 return ret
8
9 prob_sqfree_theoretical([2], [3])
```

```
10 # 0.45594532639052
```
I assume we use the first *ℓ* primes and reuse some of the same logic as above.

```
1 def density_theoretical(ps, qs):
```

```
2 ret = 1
3 for p in ps:
4 ret *=(1/p)5 for q in qs:
6 ret *= (q-1)/q7 return ret
8
9 def strategy(ell):
10 correct = 0
11 for n in range(2**ell):
12 ps, qs = binary_to_prime_sets(n, length=ell)
13 psqfree = prob_sqfree_theoretical(ps, qs)
14 density = density_theoretical(ps, qs)
15 prob = max(psqfree, 1 - psqfree)
16 correct += density * prob
17 return correct
18
19 strategy(1)20 # 0.7026423672846756
21
22 strategy(10)23 # 0.7034137933079656
24
```

```
25 strategy(20)26 # 0.7034211847385363
27
28 strategy(25)
```
²⁹ # 0.7034221516869834

Computing this for anything much larger would be prohibitively computationally expensive. Without more sophisticated thinking, it seems we've hit a wall. Presumably this continues to grow, but perhaps is strictly bounded.

I pose this as an open question. It's not clear to me how hard it is to answer it.

Question 1. *What is the limiting behavior of this strategy? Can it be shown to be less than* 71 *percent?*

5. Extended remark on the utility of ML for pure mathematics

Machine learning operates as a block box. There may be many problems where it can achieve impressive results, but it might give no clues as to how it actually does its predictions.

But one place where it is useful is acting as a **one-sided oracle** to determine whether the inputs are enough to correctly evaluate an output. For example, if I wanted to determine if a particular set of inputs are sufficient to determine some behavior, I might try to feed these inputs along with known outcomes into a machine learning blackbox. If the ML soup acts with high accuracy using only those inputs, it seems more likely that those variables are indeed significant.

It's "one-sided" because the model might simply fail to model the function well. Failing to obtain high accuracy could reflect merely that the model wasn't strong enough, or there wasn't enough data, or any of a variety of points of failure that are independent of the underlying mathematical question.

And it's an "oracle" because there are no explanations for the insight. We can not ask for the workings behind the curtain.

With regard to the question above, I think I've tried such a variety of models and structures and learning rates that I suspect that knowing *n* mod *p* for the first 100 primes isn't enough to guess $\mu(n)^2$ on random input n more than 75 percent of the time. And this belief is bolstered by the failure of the ML to do better. (I've tried neural networks of various forms too, even though I haven't described those here).

But unfortunately that's not the direction the oracle sees and I don't believe in the strength of ML enough to make a conjecture or to draw a line in the sand.

Appendix A. Parsing Logs

```
1 # path and env name : THIS IS YOUR DUMP PATH
2 path = "/scratch/"
3
4 # THE EXPERIMENTS YOU WANT TO PROBE AND THE ACCURACY INDICATOR
5 indicator = "valid_arithmetic"
6 xp_env=["dld_mu_modp_and_p_sqfree"]
7
8 # SET TO TRUE IF YOU USE BEAM SEARCH
9 has_beam=False
10
11 import os
12 import pickle
13 import matplotlib.pyplot as plt
14 import glob
15 import ast
16 from datetime import datetime
17 from tabulate import tabulate
18 import numpy as np
19 from operator import itemgetter
20
21 xp_id_filter=[]
22 xp_id_selector=[]
23 unwanted_args = ['dump_path']
24 var_args = set()
25 all_args = {}
26
27 # list experiments
28 xps = [(env, xp) for env in xp_env]29 for xp in os.listdir(path+'/'+env)
30 if (len(xp_id_selector)==0 or xp in xp_id_selector)
31 and (len(xp_id_filter)==0 or not xp in xp_id_filter)]
32 names = [path + env + '/ + xp for (env, xp) in xps]
33 print(len(names),"experiments found")
34
35 # read all args
36 pickled_xp = 0
37 for name in names:
38 pa = name+'/params.pkl'
39 if not os.path.exists(pa):
40 print("Unpickled experiment: ", name)
41 continue
42 pk = pickle.load(open(pa, 'rb'))
43 all_args.update(pk.__dict__)
```

```
16 DAVID LOWRY-DUDA LAST UPDATED: 2024.10.19
44 pickled_xp += 1
45 print(pickled_xp, "pickled experiments found")
46 print()
47
48 # find variable args
49 for name in names:
50 pa = name+'/params.pkl'
51 if not os.path.exists(pa):
52 continue
53 pk = pickle.load(open(pa, 'rb'))
54 for key,value in all_args.items():
55 if key in pk.__dict__ and value == pk.__dict__[key]:
56 continue
57 if key not in unwanted_args:
58 var_args.add(key)
59
60 print("common args")
61 for key in all_args:
62 if key not in unwanted_args and key not in var_args:
63 print(key,"=", all_args[key])
64 print()
65 print(len(var_args)," variables params out of", len(all_args))
66 print(var_args)
67
68 def vars_from_env_xp(env, xp):
69 res = {}
70 pa = path+env+'/'+xp+'/params.pkl'
71 if not os.path.exists(pa):
72 print("pickle", pa, "not found")
73 return res
74 pk = pickle.load(open(pa, 'rb'))
75 for key in var_args:
76 if key in pk.__dict__:
77 res[key] = pk.__dict__[key]
78 else:
79 res[key] = None
80 return res
81
82 def get_start_time(line):
83 parsed_line = line.split(" ")
84 dt = datetime.strptime(parsed_line[2]+''+parsed_line[3],"%m/%d/%y %H:%M:%S")
85 try:
86 idx = parsed_line.index("epoch")
87 curr_epoch = int(parsed_line[idx+1])
```

```
88 except ValueError:
89 curr_epoch = ""
90 return dt, curr_epoch
91
92 def read_xp(env, xp, indics, max_epoch=None):
93 res = {"env":env, "xp": xp, "stderr":False, "log":False,
      ,→ "error":False}
94 stderr_file = os.path.join(os.path.expanduser("~"),
      ,→ 'workdir/'+env+'/*/'+xp+'.stderr')
95 nb_stderr =len(glob.glob(stderr_file))
96 if nb stderr > 1:
97 print("duplicate stderr", env, xp)
98 return res
99 for name in glob.glob(stderr_file):
100 with open(name, 'rt') as f:
101 res["stderr"]=True
102 errlines = []
103 cuda = False
104 terminated = False
105 forced = False
106 for line in f:
107 if line.find("RuntimeError:") >= 0:
108 errlines.append(line)
109 if line.find("CUDA out of memory") >= 0:
110 cuda = True
111 if line.find("Exited with exit code 1") >=0:
112 terminated = True
113
114 if line.find("Force Terminated") >=0:
115 forced = True
116 res["forced"] = forced
117
118 res["terminated"] = terminated
119 if len(errlines) > 0:
120 res<sup>["error"]</sup> = True
121 res["runtime_errors"] = errlines
122 res["oom"] = cuda
123 if not cuda:
124 print(stderr_file,"runtime error no oom")
125
126 pa = path+env+'/'+xp+'/train.log'127 if not os.path.exists(pa):
128 return res
129 res["log"] = True
130 with open(pa, 'rt') as f:
```

```
131 series = []
132 train_loss=[]
133 for ind in indics:
134 series.append([])
135 best_val = -1.0136 best_xel = 999999999.0
137 best_epoch = -1138 epoch = -1139 \t\t val = -1140 ended = False
141 nanfound = False
142 res["curr_epoch"]=-1
143 res["train_time"]=0
144 res["eval_time"]=0
145 res["pred_nr"]=[]
146 nb_sig10 = 0
147 nb_sig15 = 0
148 counter = 0
149 counting = False
150 for line in f:
151 try:
152 if counting:
153 counter += 1
154 if line.find("Signal handler called with signal
               - 10") >= 0:
155 nb_sig10 += 1156 if line.find("Signal handler called with signal
               ,→ 15") >= 0:
157 nb_sig15 += 1
158 if line.find("Stopping criterion has been below
               \rightarrow its best value for more than") >=0:
159 ended = True
160 elif line.find("============ Starting epoch")
               \rightarrow >=0:
161 dt, curr_epoch = get_start_time(line)
162 if curr_epoch == max_epoch: break
163 res["start_time"] = dt
164 164 if curr_epoch >0 and curr_epoch ==
                  \rightarrow res["curr_epoch"]+1:
165 res["eval_time"] += (dt -
                     \rightarrow res["end_time"]).total_seconds()
166 res["curr_epoch"] = curr_epoch
167 elif line.find("============ End of epoch")
               \rightarrow >=0:
```


```
20 DAVID LOWRY-DUDA LAST UPDATED: 2024.10.19
207 res["best dic"] = dic
208 for i, indic in enumerate(indics):
209 if indicator+"_"+indic in dic:
210
                              ,→ series[i].append(dic[indicator+"_"+indic])
211
212 except Exception as e:
213 print(e, "exception in", env, xp)
214 continue
215 except:
216 print(line)
217 continue
218 res["nans"] = nanfound
219 res["ended"] = (ended or (nb_sig15 > nb_sig10))
220 res["last_epoch"] = epoch
221 res["last_acc"] = "{:.2f}".format(val)
222 res["best_epoch"] = best_epoch
223 res["best_acc"] = float("\{-.2f\}".format(best_val))
224 res["best_xeloss"] = "{:.2f}".format(best_xel)
225 res["train_loss"]=train_loss
226 res["avg_d"] = np.median(res['pred_nr'])
227 res["last_d"] = res['pred_nr'][-1] if
          \rightarrow len(res['pred_nr']) > 0 else -1
228 if epoch >=0:
229 res["train_time"] /= (epoch+1)
230 res["eval time"] / = (epoch+1)
231 res["train_time"] = int(res["train_time"]+0.5)
232 res["eval_time"] = int(res["eval_time"]+0.5)
233
234 for i,indic in enumerate(indics):
235 res['last-"+indic] = "{:.}2f}'.format(series[i][-1]),→ if len(series[i])>0 else '0'
236 res["best_"+indic] =
              \rightarrow "\{:\, 2f\}".format(max(series[i])) if
              \rightarrow len(series[i])>O else 'O'
237 res[indic] = series[i]
238 if len(series[i])!= epoch + 1:
239 print("mismatch in nr of epochs",env, xp,
                 \rightarrow epoch+1, len(series[i]), indic)
240 return res
241
242 data = []243 indics = ["beam_acc" if has_beam is True else "acc","xe_loss"]
```

```
244 indics.extend(["correct", "perfect", "beam_acc_d1",
    ,→ "beam_acc_d2",
245 "beam_acc_nb", "additional_1","additional_2","additional_3"])
246
247 for (env, xp) in xps:
248 res = read_xp(env, xp, indics, None) # USE THE LAST
        ,→ PARAMETER IF YOU WANT TO LIMIT READ TO N EPOCHS
249 res.update(vars_from_env_xp(env, xp))
250 data.append(res)
251
252 print(len(data), "experiments read")
253 print(len([d for d in data if d["stderr"] is False]),"stderr
    \rightarrow not found")
254 print(len([d for d in data if d["error"] is True]),"runtime
    ,→ errors")
255 print(len([d for d in data if "oom" in d and d["oom"] is
    ,→ True]),"oom errors")
256 print(len([d for d in data if "terminated" in d and
    \rightarrow d["terminated"] is True]), "exit code 1")
257 print(len([d for d in data if "forced" in d and d["forced"] is
    \rightarrow True]), "Force Terminated")
258 print(len([d for d in data if "last_epoch" in d and
    \rightarrow d["last_epoch"] >= 0]), "started XP")
259 print(len([d for d in data if "ended" in d and d["ended"] is
    ,→ True]),"ended XP")
260 print(len([d for d in data if "best_acc" in d and
    ,→ float(d["best_acc"]) > 0.0]),"began predicting")
```
And to make some graphs displaying various things, I would run the following. Or rather, I would run the above and below in a notebook, so the graphs display inline. (Otherwise I guess I would save them).

In practice, it was sufficient to look at the tail of the running log and to extract learning rate failures and accuracies on test sets.

```
1 import numpy as np
2
3 def compose(f,g):
4 return lambda x : f(g(x))5
6 def print_table(data, args, sort=False):
7 \text{ res} = [1]8 for d in data:
9 line = [d[v] if v in d else None for v in args]
10 res.append(line)
11 if sort:
```

```
22 DAVID LOWRY-DUDA LAST UPDATED: 2024.10.19
12 res = sorted(res, key=compose(float,itemgetter(0)),
          ,→ reverse=True)
13 print(tabulate(res,headers=args,tablefmt="pretty"))
14
15 def speed_table(data, args, indic, sort=False, percent=95):
16 res = \lceil \rceil17 for d in data:
18 if indic in d:
19 line = [d[v] if v in d else None for v in args]
20 val= 10000
21 for i, v in enumerate(d[indic]):
22 if v > = percent and i < val:
23 val = i
24 line.insert(1,val)
25 res.append(line)
26 e= \arg s \cdot \text{copy}()27 e.insert(1,'first epoch')
28 if sort:
29 res = sorted(res, key=compose(float,itemgetter(1)),
          ,→ reverse=False)
30 print(tabulate(res,headers=e,tablefmt="pretty"))
31
32 def training_curve(data, indic, beg=0, end=-1, maxval=None,
   ,→ minval=None, export_to=""):
33 print(indic)
34 for d in data:
35 if indic in d:
36 if end == -1:
37 plt.plot(d[indic][beg:],linewidth=1)
38 else:
39 plt.plot(d[indic][beg:end],linewidth=1)
40 plt.ylim(minval,maxval)
41 plt.rcParams['figure.figsize'] = [10,10]
42 if export_to != '':
43 # print(export_to)
44 plt.savefig(export_to,bbox_inches="tight")
45 plt.show()
46
47 def filter_xp(xp, filt):
48 for f in filt:
49 if not f in xp:
50 return False
51 if not xp[f] in filt[f]:
52 return False
```

```
53 return True
54
55 def xp_stats(data, splits, best_arg, best_value):
56 res dic = \{ \}57 nb = 0
58 for d in data:
59 if d[best_arg] < best_value: continue
60 nb += 1
61 for s in splits:
62 if not s in d: continue
63 lib=s+':'+str(d[s])
64 if lib in res_dic:
65 res_dic[lib] += 1
66 else:
67 res dic[lib] = 1
68 print()
69 print(f"{nb} experiments with accuracy over {best_value}")
70 for elem in sorted(res_dic):
71 print(elem,' : ',res_dic[elem])
72 print()
73
74 xp_filter ={}
75
76 # CHANGE THESE TO FILTER THE EXPERIMENTS
77 #xp_filter.update({"n_enc_layers":[4]})
78 #xp_filter.update({"enc_emb_dim":[512]})
79
80 fdata = [d for d in data if filter_xp(d, xp_filter) is True]
81
82 oomtab = [d for d in fdata if d['error"] is True]
83 print(f"CUDA out of memory ({len(oomtab)})")
84 print_table(oomtab, var_args)
85
86 forcetab = [d for d in fdata if 'forced' in d and d["forced"]
   \rightarrow is Truel
87 print(f"Forced terminations ({len(forcetab)})")
88 print_table(forcetab, var_args)
89
90 unstartedtab = [d for d in fdata if "last_epoch" in d and
   \rightarrow d["last_epoch"] < 0]
91 print(f"Not started ({len(unstartedtab)})")
92 print_table(unstartedtab, var_args)
93
94 crypto = False
```

```
24 DAVID LOWRY-DUDA LAST UPDATED: 2024.10.19
95 runargs = ["best_acc", "best_epoch","best_xeloss", "ended",
    ,→ "last_epoch",
96 "last_acc", "last_xe_loss","nans", "error", "train_time",
    ightharpoonup "eval time"]
97
98 #runargs.extend(["best_acc_d1" , "best_acc_d2"])
99 for v in var_args:
100 runargs.append(v)
101 runningtab = [d for d in fdata if "last_epoch" in d and
    \rightarrow d["last_epoch"] >= 0]
102 print(f"Running experiments ({len(runningtab)})")
103
104 #splits = ['n_enc_layers','dec_emb_dim','reload_size']
105 #xp_stats(fdata, splits, 'best_acc',90.0)
106 print()
107 print_table(runningtab, runargs, sort=True)
108
109 training_curve(fdata, "beam_acc" if has_beam is True else
    \rightarrow "acc", 0, -1, None, export_to = "")
110 training_curve(fdata, "perfect")
111 training_curve(fdata, "correct")
112
113 training_curve(fdata, "xe_loss", 0) #, None, 0.9* np.min([x for\rightarrow d in fdata for x in d["xe_loss"] if x > 0.0])
114 training_curve(fdata, "train_loss",0, -1, 2)
115 speed_table(runningtab, runargs, "beam_acc" if has_beam else
    ,→ "acc", sort=True,percent=99)
116 speed_table(runningtab, runargs, "beam_acc" if has_beam else
    ,→ "acc", sort=True,percent=50)
117 speed_table(runningtab, runargs, "beam_acc" if has_beam else
    ,→ "acc", sort=True,percent=55)
118 speed_table(runningtab, runargs, "beam_acc" if has_beam else
    ,→ "acc", sort=True,percent=60)
```
REFERENCES

[DLD-General] David Lowry-Duda, *General Report on Machine Learning Experiments for the Möbius Function*. 2024 October 21. (Cited on page [1\)](#page-0-1) [Int2Int] *Int2Int* Github Repository, [https://github.com/f-charton/](https://github.com/f-charton/Int2Int) [Int2Int](https://github.com/f-charton/Int2Int). Accesssed 2024 October 20. (Cited on page [1\)](#page-0-1)